



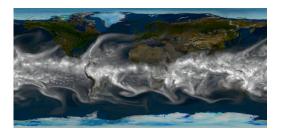
On the Role of Indexing for Big Data in Scientific Domains

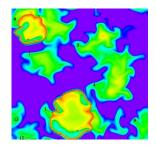
Arie Shoshani Lawrence Berkeley National Lab

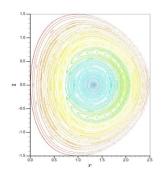
BIGDATA and EXTREME-SCALE COMPUTING April 30-May 1, 2013

Outline

- Examples of indexing needs in scientific domains
- Scientific Indexing requirements
- Bitmaps indexing as a promising technology







Example of Big Data in Science

Large Hadron Collider: to find the God particle

- **I5 PB** per year sensors capable of I40PB/s
- 27 km tunnel
- ~10,000 superconducting magnets
- Operating temperature 1.9 Kelvin
- Construction cost: US\$9Billion
- Power consumption: ~120 MW



Typical Event Figures

Experiment	# members /institutions	Date of first data	# events/ year	volume/year- TB
STAR	350/35	2001	10 ⁸ -10 ⁹	500
PHENIX	350/35	2001	10 ⁹	600
BABAR	300/30	1999	10 ⁹	80
CLAS	200/40	1997	10 ¹⁰	300
ATLAS	1200/140	2008	10 ¹⁰	5000



STAR: Solenoidal Tracker At RHIC RHIC: Relativistic Heavy Ion Collider

LHC: Large Hadron Collider Includes: ATLAS, CMS, ...

A mockup of An "event"



What are the indexing challenges?

- Generate large amounts of raw data referred to as "events"
 - ♦ Collected from simulations and experiments

Post-processing of data

- ♦ Identify elements in data (find particles produced, tracks)
- \diamond generate summary variables per event
 - \circ e.g. momentum, no. of pions, transverse energy
 - Number of variables is large (50-100)
- Analyze data
 - \diamond use summary variables to characterize events
 - \diamond extract subsets from the large dataset
 - Need to access events based on partial variable specification (range queries)
 - \circ e.g. ((0.1 < AVpT < 0.2) ^ (10 < Np < 20)) v (N > 6000)

Challenges

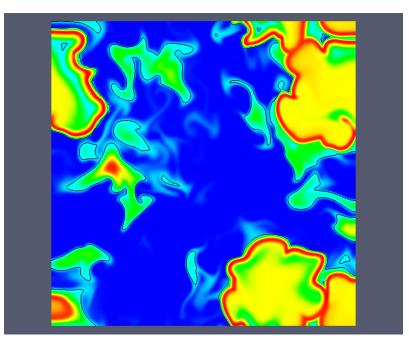
- \diamond Search over billions of events
- \diamond Multi-variable search, but only over a subset of the variable
- ♦ Type of query: a needle-in-the-haystack
- \diamond Another type of query: larger subsets for statistical properties
- ♦ Search over numerical values (integers, floating point)

Combustion simulation example

- Combustion simulation: 1000x1000x1000 mesh with 100s of chemical species over 1000s of time steps – 10¹⁴ data values
- This is an image of a <u>single</u> variable (temperature)
- What's needed is search over multiple variables, such as: Temperature > 1000 AND pressure > 106 AND HO2 > 10⁻⁷ AND HO2 > 10⁻⁶

Challenges

- Multi-variable queries from a subset of variables
- □ Search over numerical values
- Identify large number of regions



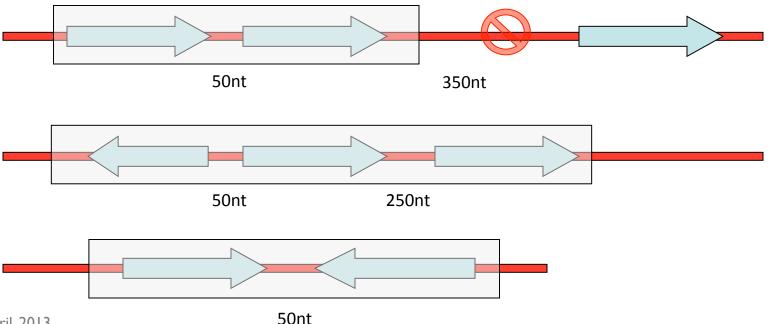
Gene functional annotation

Sequence similarity

Gene context provide information for the function of genes.

Functionally related genes are frequently found in the same chromosomal neighborhood.

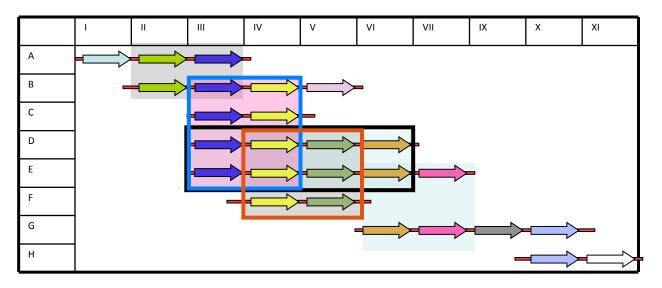
- Cassette Definition
 - Parallel or Divergent orientation
 - Distance < 300nt



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Conserved chromosomal cassettes

- Genes are replaced by protein families (COGs, pfams, IMG ortholog families). One gene → multiple families.
- We refer to these as "properties", such as "cog0087 cog0088 cog0089 pfam00181 pfam00189 pfam00203 pfam00237"



- Boxes that share two cassettes and two genes, if the genomes are distant phylogenetically (more than species)
- E.g. for black box: blue and red are 1st step relatives.

Why is this problem hard?

🖵 Size

I00 million cassettes, with properties from about 25,000 possible values (currently).

 \diamond Total number of elements: 2.5 x 10¹²

Challenge

Query types

Given a cassette find all cassettes that have the same properties in common

 $\,\circ\,$ That is a massive multi-value search

 Given a cassette find all cassettes that have 2-or-more properties in common (in general k-or-more)

 \circ Explosive search of all possible combinations of 2-or-more

Big Data Indexing Requirements

□ Speed of search

 \diamond Search over billions – trillions data values in seconds

Multi-variable queries

 \diamond Be efficient for combining results from individual variable search results

□ Size of index

 \diamond Index size should be a fraction of original data

Granularity

Ability to produce smaller indexes when granularity can be reduced, such as I decimal points, for example

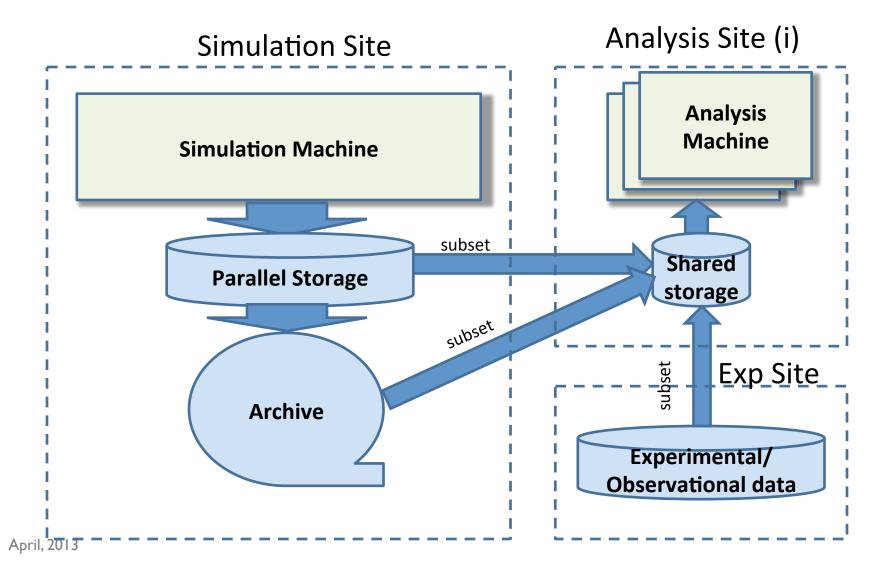
Parallelism

 \diamond Should be easily partitioned into sections for parallel processing

□ Speed of index generation

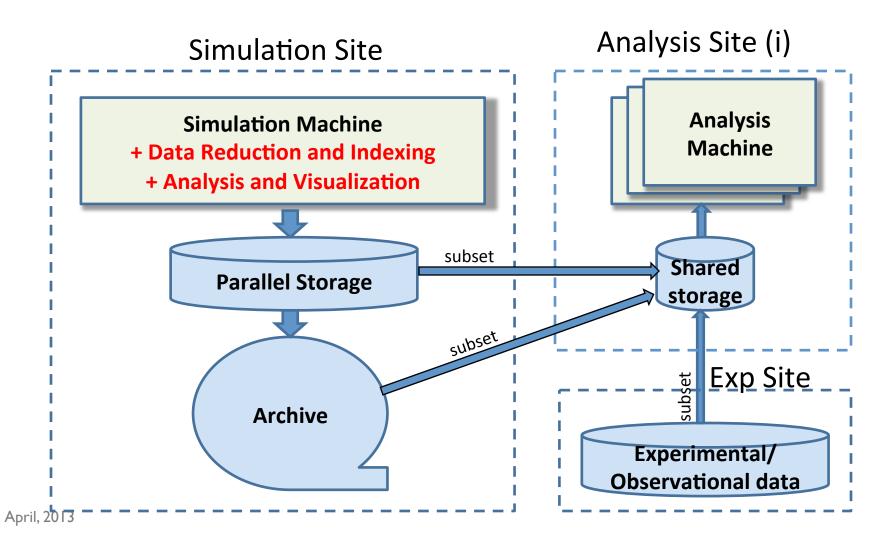
 \diamond For in situ processing, index should be built at the rate of data generation

Scaling simulations generates a data volume challenge (PBs)



What Can be Done?

Perform some data analysis and visualization on simulation machine (in-situ)
 Reduce Data and prepare data for further analysis (in-situ)



Data Analysis

Two fundamental aspects

- Pattern matching: Perform analysis tasks for finding known or expected patterns
- Pattern discovery: Iterative exploratory analysis processes of looking for unknown patterns or features in the data

Ideas for the analysis of Big Data

- \diamond Perform pattern matching tasks in the simulation machine
 - \circ "In situ" analysis
- Prepare data for pattern discovery on the simulation machine, and perform analysis on mid-size analysis machine
 - \circ "In-transit" data preparation
 - \circ "Off-line" data analysis

Index:

A Data Structure for Accelerating Data Accesses

Tree-based indexes

- \diamond E.g. family of B-Trees
- \diamond Commonly used database management systems
- \diamond Sacrifice search efficiency to permit dynamic update

Multi-dimensional indexes

- \diamond E.g. R-tree, Quad-trees, KD-trees, ...
- \diamond Don't scale for large number of dimensions
- \diamond Are inefficient for partial searches (subset of attributes)

Hashing

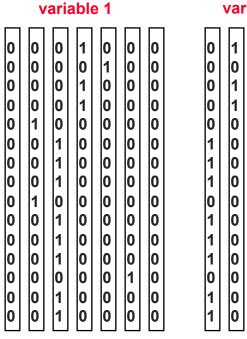
- ♦ Predictable performance
- \diamond Good for locating individual data records

Bitmap indexes:

- \diamond Good for read-mostly data
- \diamond Handle partial range queries efficiently
- \diamond May have trouble handling data with a large number of distinct values

Bitmap index for each variable

- Take advantage that index need to be is append only
- Generate a bitmap for each possible value of each variable
 - (e.g. for 0<Np<300, have 300 bitmaps)
- compress each bit vector (some version of run length encoding)
- Need to touch only bitmaps for the specified search



variable 2

0 0

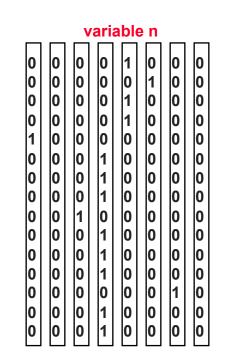
0 0 0

0 0

0 0

0 0

. .



Basic Bitmap Index

Data $b_0 b_1 b_2 b_3 b_4 b_5$ =0 =1 =2 =3 =4 =5 values О 5 3 2

- Easy to build: faster than building B-trees
- <u>Efficient for querying</u>: only bitwise logical operations
 - $A < 2 \rightarrow b_0 \text{ OR } b_1$
 - $A > 2 \rightarrow b_3 OR b_4 OR b_5$
- Efficient for multi-dimensional queries
 - Use bitwise operations to combine the partial results
- Size: one bit per distinct value per object
 - Definition: Cardinality == number of distinct values
 - Need to control size for high cardinality attributes
- Main idea:
 - highly efficient compression method
 - Compute friendly can perform operations directly on compressed data

FastBit properties – highly efficient and compact

Main idea:

 \diamond Invented specialized compression methods (was patented) that:

- Can perform logical operations directly on compressed bitmaps
- Excels in support of multi-variable queries
- Can partition and merge bitmaps without decompression essential for parallelization of indexes

FastBit takes advantage of append only data to achieve:

- Search speed by I0x I00x than best known bitmap indexing methods
- On average about 1/3 of data volume compared to 2-3 times in common indexes because of compression method
- Proven to be theoretically optimal data search time is proportional to size

of the result

Usage

- ♦ In multiple scientific application in DOE
- Embedded into in situ frameworks
- Thousands of downloads around the world (open source under source forge), including commercial companies

Methods to Improve Bitmap Index

Compression

- ♦ FastBit compression method: Word-Aligned Hybrid (WAH) code
- ♦ I0x speedup over Byte-aligned Bitmap Code

Encoding

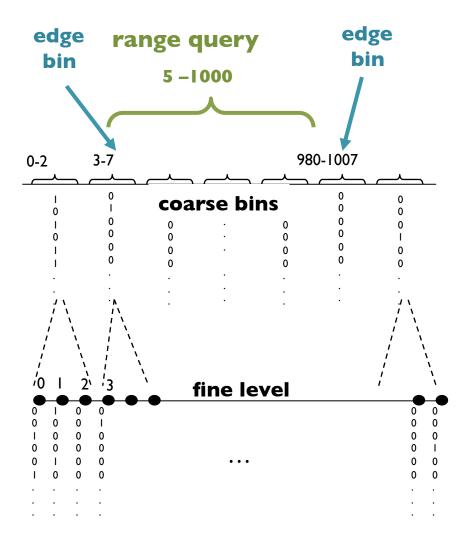
- \diamond Multi-level encoding
- \diamond Reduce bitmaps needed for a query
- \diamond 5x speedup

Binning

- \diamond Some times we choose to use bins at the fine level to reduce index size
- \diamond Problem: if query falls in the middle of edge bins
- Solution: Order-preserving Bin-based Clustering (OrBiC)
- \diamond 5x speedup for searching bins

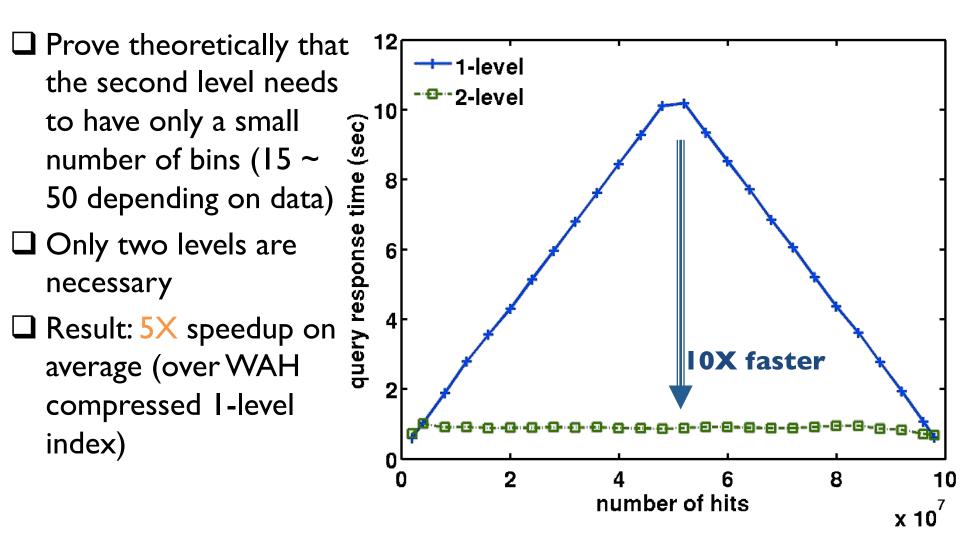
Improving Bitmap Indexes: Multi-Level Encoding

- The finest-level may be precise or binned
- Coarse levels are always binned
- Each coarse bin contains a number of fine bins/values
- Queries can be processed with a combination of coarse and fine bitmaps
- Only edge bins need to be resolved at the fine level
- Analysis revealed how to construct the coarse level in order to reduce the query processing time



[Wu, Shoshani and Stockinger 2010]

Two Levels Are Better Than One



Domain-Specific Challenges – current and future

Generate index at the rate of data generation in situ

 \diamond Increase level of parallel processing

 \diamond Perform partial index generation per node

 \diamond Take advantage of local NVRAM

□ Adapt indexing methods to a variety of data models

 \diamond Irregular grids, geodesic meshes, toroidal meshes

 $\circ\,$ How to linearize the space

Multi-level grid, such as adaptive-mesh-refinement (AMR)

Use results of index in subsequent operations

 \diamond Statistical summaries

 \diamond Region growing, region overlaps, ...

□ Adapt indexing to specialized operations

- \diamond Searches for k-or-more matches
- ♦ Searches based on formulas (plug-in-codes)

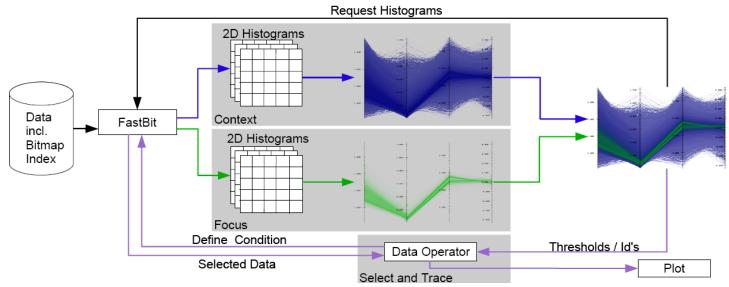
FastBit in support for Query-Driven Visualization

Collaboration between SDM and Vis groups

• Use FastBit indexes to efficiently select the most interesting data for visualization

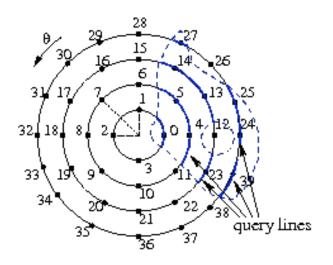
Example: laser wakefield accelerator simulation

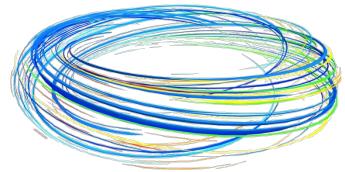
- VORPAL produces 2D and 3D simulations of particles in laser wakefield
- Finding and tracking particles with large momentum is key to design the accelerator
- Brute-force algorithm is quadratic (taking 5 minutes on 0.5 mil particles), FastBit time is linear in the number of results (takes 0.3 s, 1000 X speedup)



FastBit adaptation to toroidal meshes

- Extended FastBit indexing capability to search for regions of interest defined on toroidal meshes used for fusion simulations
- Developed algorithms to take full advantages of the regularity present in the magnetic coordinates but not in the Cartesian coordinates
- Much more compact than the general connectivity graph: ~ 200 numbers vs. 6 million numbers
- Labeling query lines using magnetic coordinates is 600-1000 x faster than using connectivity graph
- Developed new Connected Component Labeling
 algorithm
 - Recently used in Atmospheric Rivers project





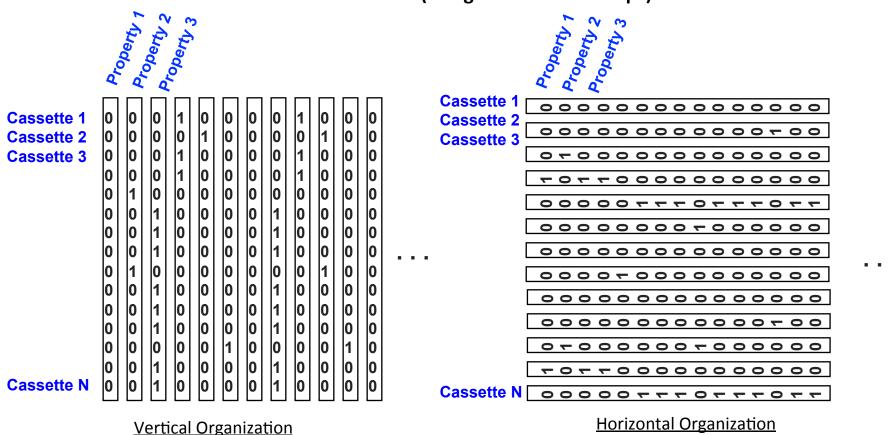
Adapting to cassette searches

Results: (1) given a cassette, search all cassettes with the same properties

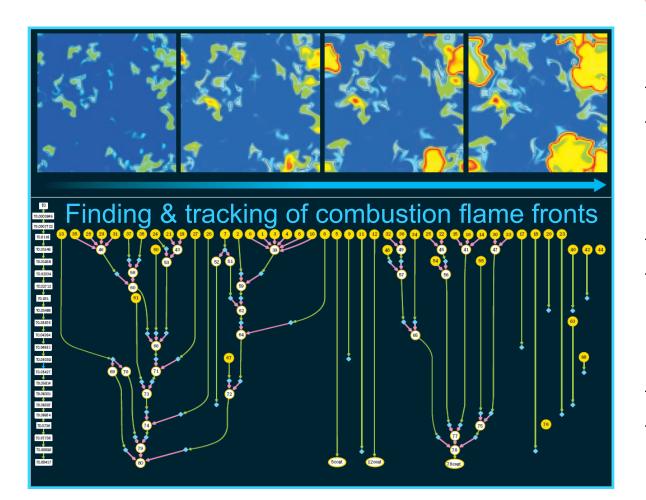
Done in about 0.07 second (using vertical bitmaps)

(2) Find similar cassettes with 2 or more properties

Done in about 10-15 seconds (using horizontal bitmaps)



Flame Front Tracking in Combustion



Challenges

♦ Cell identification

☆Identify all cells that satisfy range conditions

♦ Region growing
 ♦ Connect neighboring cells into regions

♦ Region tracking

Big Data Indexing Requirements: Bitmap indexing advantage

□ Speed of search

- ♦ Search over billions trillions data values in seconds
 - Yes, with compute-friendly compression

Multi-variable queries

- \diamond Be efficient for combining results from individual variable search results
 - $\circ~$ Yes, combining results for each variable as bitmaps is very efficient

Size of index

- \diamond Index size should be a fraction of original data
 - Yes, compression of bitmap index is essential

Granularity

- Ability to produce smaller indexes when granularity can be reduced, such as 2 decimal points, for example
 - $\circ~$ Yes, binning over multiple values proved very effective

] Parallelism

- \diamond Should be easily partitioned into sections for parallel processing
 - Yes, if compressed bitmaps can be easily combined (WAH has this property)

Speed of index generation

- $\diamond\,$ For in situ processing, index should be built at the rate of data generation
 - OK for billions of values, but a trillion value index took 10 minutes (still a challenge)

Architectural Changes that Could Benefit in situ indexing

NVRAM on each node

- \diamond Can be used to build partial indexes over multiple time steps
- \diamond Can be used to accelerate in-situ index generation

□ Take advantage of GPUs

 \diamond Assign index generation for each variable to separate GPUs

□ NVRAM between machine and storage system

Can be used for generating indexing for post-processing while data is streaming out to be stored on disk





THANKS!

Co-authors:

E.W. Bethel, S. Byna, J. Chou, W. S. Daughton, M. Howison, H. Karimabadi, K.-J. Hsu, K.-W. Lin, V. Markowitz, K. Mavrommatis, Prabhat, A. Romosan, V. Roytershteynz, O. Rübel, A. Shoshani, A. Uselton

FastBit software <u>http://sdm.lbl.gov/fastbit/</u> FastQuery software <u>http://goo.gl/iBw6V</u> Scientific Data Management group <u>http://sdm.lbl.gov/</u>



